Constraint Reasoning and Motion Planning in the SAUNA Project

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<u>L</u>Motivation

AGVs for Industrial Automation

Autonomous Ground Vehicles (AGVs) becoming paramount to industrial automation

mining (e.g., Atlas-Copco) construction (e.g., Volvo) logistics (e.g., Kollmorgen)

- Several key processes are still ad-hoc and labor-intensive
	- AGV paths often pre-defined and hand-crafted
	- crude planning/allocation/scheduling heuristics
	- conflicts "avoided off-line" rather than resolved on-line
	- lack of **flexibility** and **reconfigurability**

LMotivation

What If.

- Site operators could **post high-level requirements**
	- new tasks, vehicles *e.g., "pick up new incoming loads"*
	- spatial constraints *e.g., "zone A is now off-limits"*
	- temporal constraints *e.g., "complete mission of vehicle A within 10 minutes"*
- Vehicles could **adapt** their current and scheduled trajectories accordingly
	- vehicles coordinate automatically in response to posted requirements
- **Performace of vehicles** could be automatically taken into account by **other vehicles**

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Motivation

The Core Issue of Coordination

- Coordination is necessary to ensure absence of **collisions** and **deadlocks**
- **Current practice:** compute trajectories for all vehicles **before** coordination occurs
	- collisions and deadlocks are dealt with locally
	- o overall fleet requirements cannot be guaranteed
- **Our view:** all decisions regarding vehicle trajectories can be seen as **constraints on trajectories**
	- some constraints **known in advance**, some need to be **inferred**
- **Least commitment:** impose increasigly tight constraints on trajectory, but **do not commit** to specific solution until **execution**

LMotivation

2 [Conflict Resolution as Refinement of Trajectory Envelopes](#page-14-0)

3 [Control with Trajectory Envelopes](#page-18-0)

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 $\mathcal{A} \cong \mathcal{B} \times \mathcal{A} \cong \mathcal{B}$

Non-Holonomic Vehicles

- Vehicles are non-holonomic \Rightarrow no sideways motion
- "Car-like" vehicles

Kinematic Model

$$
\dot{q} = f(q, v) = (v \cos(\theta), v \sin(\theta), \frac{v}{l} \tan(\phi), \dot{\theta})
$$

$$
q = (x, y, \theta, \phi) \in \mathbb{R}^4 \text{ (state vector)}
$$

$$
v = (v, \dot{\theta}) \in \mathbb{R}^2 \text{ (control vector)}
$$

Path: $p: [0,1] \to \mathbb{R}^2$ (parametrized using arc length)

• Trajectory: $p(\sigma(t))$ (σ = time history along path)

Trajectories and Trajectory Envelopes (I)

- **Trajectory envelope: spatial** requirements on *p* and **temporal** requirements on $\sigma(t)$
	- **temporal envelope**, **spatial envelope**

Trajectories and Trajectory Envelopes (II)

- **•** Trajectory $p(\sigma(t))$ is **admissible** if
	- it can be obtained from the evolution of the **kinematic model**
	- \bullet *p* / $\sigma(t)$ lies within the **spatial** / **temporal** envelope

Explicit Requirements as Constraints

- Avoiding **fixed obstacles** = imposing **spatial constraits** on trajectory envelopes
	- drivable area encoded as spatial constraints
	- any additional spatial requirements can be posted as well
- High-level **temporal requirements** = imposing **temporal constraints** on trajectory envelopes
	- deadlines, durations, and any other temporal requirements can be posted directly on trajectory envelopes
- We can constrain **admissible trajectories** to account for all high-level spatial and temporal requirements
	- however, **collisions** between vehicles and **deadlocks** are still possible. . .

Computing Spatial Envelopes (I)

- Paths computed over drivable area using *ARA*[∗]
	- anytime path planning algorithm [\[Likhachev et al., 2003\]](#page-31-1)
	- uses kinematically feasible motion primitives

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Trajectory Envelopes

Computing Spatial Envelopes (II)

- Spatial envelopes computed over obtained paths
	- sample paths with given $\Delta\sigma \propto (1/c$ urvature of the path)
	- calculate polyhedron $\mathscr{S}^{(j)}_i$ i ^{ov} enclosing the sampled point

Computing Temporal Envelopes

Impose temporal bounds on **spatial envelope traversal**

$$
\ell_i^{(j)} \le e_i^{(j)} - s_i^{(j)} \le u_i^{(j)}
$$

$$
\ell_{i,i+1}^{(j)} \le e_i^{(j)} - s_{i+1}^{(j)} \le u_{i,i+1}^{(j)}
$$

Bounds computed with **slowest** and **fastest** speed traversal of reference path

$$
\ell_i = e_i^{\text{fast}} - s_i^{\text{slow}}, \quad u_i = e_i^{\text{slow}} - s_i^{\text{fast}},
$$

$$
\ell_{i,i+1} = e_i^{\text{fast}} - s_{i+1}^{\text{slow}}, \quad u_{i,i+1} = e_i^{\text{slow}} - s_{i+1}^{\text{fast}}
$$

L Conflict Resolution as Refinement of Trajectory Envelopes

Outline

2 [Conflict Resolution as Refinement of Trajectory Envelopes](#page-14-0)

[Control with Trajectory Envelopes](#page-18-0)

[Evaluation](#page-24-0)

 $\left\{ \begin{array}{ccc} 1 & 0 & 0 \\ 0 & 1 & 0 \end{array} \right\}$, $\left\{ \begin{array}{ccc} \frac{1}{2} & 0 & 0 \\ 0 & 0 & 0 \end{array} \right\}$, $\left\{ \begin{array}{ccc} \frac{1}{2} & 0 & 0 \\ 0 & 0 & 0 \end{array} \right\}$

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Conflict Resolution as Refinement of Trajectory Envelopes

Spatio-Temporal Conflicts

- Polygons that overlap in **time** and **space** constitute a **conflict** ۰
- **Eliminating the temporal overlap** resolves a conflict 0

Spatio-temporal overlap:

$$
\begin{array}{c} \mathscr{S}^{(j)}_i \cap \mathscr{S}^{(m)}_k \neq \mathbf{0} \wedge \\ \left[s^{(j)}_i, e^{(j)}_i \right] \cap \left[s^{(m)}_k, e^{(m)}_k \right] \neq \mathbf{0} \end{array}
$$

Semantics of **temporal** overlap:

 $\left[s_i^{(j)}, e_i^{(j)}\right] \cap \left[s_k^{(m)}, e_k^{(m)}\right] \neq \emptyset$ if $\max\left(\underline{s}_i^{(j)}, \underline{s}_k^{(m)}\right) \leq \min\left(\underline{e}_i^{(j)}, \underline{e}_k^{(m)}\right)$

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Conflict Resolution as Refinement of Trajectory Envelopes

Conflict Resolution as a Meta-CSP (I)

- Spatio-temporal conflicts solved by **posting temporal constraints** [\[Cesta et al., 2002\]](#page-31-2)
	- remove temporal overlap between polygons
- Search for resolving constraints **cast as a CSP**
	- **variables:** pairs of spatio-temporally overalpping polygons
	- **values:** temporal constraints that eleminate temporal overlap
- Collisions are avoided by **refining the trajectory envelopes**
	- adding temporal constraints to eliminate temporal overlap
- **Note:** temporal envelope refinement can be alternated with **spatial envelope refinement**
	- adding spatial constraints to eliminate spatial overlap

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Conflict Resolution as Refinement of Trajectory Envelopes

Conflict Resolution as a Meta-CSP (II)

- Search for resolving temporal constraints performed using standard backtracking search algorithm
- **Variable ordering heuristic** based on **spatial** features
	- prefer pairs of polygons that are **spatially closer** to other conflicting pairs
- **Value ordering heuristic** based on **temporal** features
	- prefer orderings that least impact overall **temporal flexibility**
- Conflict resolution also eliminates **deadlocks**
	- backtracking search breaks cycles on resource usage (resources = intersections of temporally overlapping polygons)

Outline

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From Trajectory Envelopes to Control Actions (I)

- Conflict resolution **refines the temporal envelope** to obtain conflict- and deadlock-free trajectory envelopes
- We now have available **many alternative trajectories** that can be sampled from these envelopes
	- spatial constraints designate the polygons in which vehicles should be
	- **•** temporal constraints determine flexible bounds on each polygon

From Trajectory Envelopes to Control Actions (II)

- Each vehicle has a **tracking controller**
	- computes **kinematically feasible** control actions for vehicle
	- minimizing divergence from a **reference trajectory**
	- while accounting for **spatial constraints**
	- based on efficient QP solver [\[Dimitrov et al., 2011\]](#page-31-3)
- **Controller input** = **nominal path** + **fixed temporal bounds** + **spatial envelope**

Sampling Trajectory Envelopes (I)

- Due to efficient formulation, controller can compute control actions for **several alternative** reference trajectories
	- **choose** "best" one to follow according to **trakcing performance**
- Alternative bounds can be computed in **polynomial time**
	- controllers choose current alternative **consistently** w/ other vehicles (executed temporal profiles must be solution of the STP)

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L Fvaluation

Quantitative Evaluation on Artificial Problems

- \bullet 900 problems
- 9 test sets (2–10 vehicles) 0
- **•** randomly generated paths
- 100 trials per test set 0
- Note: artificially difficult 0 problems

L Evaluation

Quantitative Evaluation on Artificial Problems

- Average conflict resolution time $<$ 1 sec with up to 8 vehicles \bullet
- 0 Maximum time to calculate an alternative temporal profile $<$ 50 msec

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L Fvaluation

Realistic Test Run

- 7 Linde vehicles, 140 polyhedra, pre-assigned start/destination poses
- Path planning: 5 sec / Conflict resolution: 34 sec (13 temporal constraints) 0

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Time to calculate alternative temporal profiles: 250 msec 0

L Conclusions

Conclusions

- Fleet management typically consists of **several processes**
	- all processes impose requirements on trajectories
- **Constraint-based formulation** provides processes with a **uniform model**
	- perception, planning, scheduling, control all contribute to refining trajectory envelopes
- Not all proccesses are necessarily **automated**
	- \bullet human operators, a pre-existing dispatching strategy, \dots
	- this is important for industrial application
- Modular approach maximizes the ability to **react at different levels**
	- modules impose increasingly low-level constraints on overall solution
	- commitment to specific trajectories done **as late as possible**

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$\mathsf{\mathsf{L}}$ Conclusions

SAUNA Functional View

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